REMARKS

Entry of the foregoing amendment is respectfully requested. This amendment is intended to place the claims in a more conventional format and eliminate the multiple dependency of the claims.

Respectfully submitted,

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Attachment to Preliminary Amendment dated September 10, 2001

Marked-up Claims 1-7

1.   (Amended) A countermeasure method against attacks by differential analysis in an electronic component implementing a secret key [(K)] cryptographic algorithm, the implementation of which comprises a number of successive calculation cycles [(T1, ... T16)] in order to supply, from first input data [(L0, R0)] applied to the first cycle [(T1)], final data [(L16, R16)] at the output of the last cycle [(T16) allowing the production of] to produce an encrypted message [(C)], each calculation cycle using calculation means [(TC)] for supplying an output data item [(S)] from an input data item [(E)], said calculation means [comprising the application of] performing the steps of:

applying a first random value [(u)] to the input data item [(E)] and to the output data item [(S)] in order to obtain [at the output] an unpredictable data item [(S⊕u), characterised in that the method comprises the use of means of] as an output, and

applying a second random value [(v)] to said first input data [(L0, R0), according to] by means of an EXCLUSIVE OR operation.

2.   (Amended) A countermeasure method according to Claim 1, [characterised in that it also comprises the use of means] further including the step of applying the second random value [(v)] to the final data supplied by the last cycle [(T16), according to] by means of an EXCLUSIVE OR operation.

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3.   (Amended) A countermeasure method according to [either one of the previous claims, characterised in that it comprises] claim 1 further including the step, at the end of each cycle, [the execution] of executing an additional operation [(CP(p(u))) in order] to eliminate said first random value [(u)] at the output of each cycle.

4.   (Amended) A countermeasure method according to [any one of the previous claims, characterised in that it comprises the taking of] claim 1 wherein a new set of first and second random values [(u, v) and calculation of the calculation means (TC_M) used in each cycle] is selected for each new execution of the algorithm.

5.   (Amended) A method according to Claim 4, [characterised in that] wherein said calculation means [(TC_M)] are calculated from first calculation means [(TC_0)] defining, for input data [(E)], corresponding output data [(S)], by applying the second random value [(v)] to said input data [($E \oplus e(v)$)] and applying the first random value [(u)] at least to said output data [($S \oplus u$)] of the first calculation means.

6.   (Amended) A countermeasure method according to Claim 5, [characterised in that] wherein the calculation means [(TC_0, TC_M) are] comprise constants tables.

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7.   (Amended) An electronic security component [implementing the] <u>that</u>

<u>implements a</u> countermeasure method [against attacks by differential analysis comprising]

<u>for attacks against</u> a secret key [(K)] cryptographic algorithm[, the implementation of

which] <u>by means of differential analysis, wherein said algorithm</u> comprises a number of

successive calculation cycles [(T1, ... T16)] in order to supply, from first input data [(L0,

R0)] applied to the first cycle [(T1)], final data [(L16, R16)] at the output of the last cycle

[(T16) allowing the production of] <u>to produce</u> an encrypted message [(C)], each calculation

cycle using calculation means [(TC)] for supplying an output data item [(S)] from an input

data item [(E)], said calculation means comprising the application of a first random value

[(u)] to the input data item [(E)] and to the output data item [(S in order] to obtain [at the

output] an unpredictable <u>output</u> data item [(S⊕u), characterised in that]<u>, comprising</u> first

calculation means [(TC$_0$) are] fixed in <u>a</u> program memory [(1) of said component], <u>second</u>

calculation means [(TC$_M$) used in each cycle being] <u>that are</u> calculated at each new

execution of the algorithm and stored in working memory [(3)], and [in that it comprises]

means [(4) of] <u>for</u> generating first and second random values [(u, v)] for calculating said

<u>second</u> calculation means [(TC$_M$)].